## Understanding Information Disclosure from Secure Computation Output

Analytical and Data-Driven Analysis

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April 16, 2025

#### Motivation for secure computation





 $f(s_1, s_2, s_3, s_4, s_5) = o$ 







 Want to perform computation on private data

#### Motivation for secure computation



- Want to perform computation on private data
- Employ cryptographic techniques to compute without "seeing" the data

- Emphasis has been on the techniques
- Efficiency, performance improvements
- Development tools to improve accessibility



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## Is this sufficient?

**Technical Report** 

CORPORATION

## Achieving Higher-Fidelity Conjunction Analyses Using Cryptography to Improve Information Sharing

Brett Hemenway, William Welser IV, Dave Baiocchi

"There are many situations, however, where the output of the protocol itself may leak too much information... this leakage seems to be acceptable to the community, but this is a question that needs to be addressed before any MPC protocol can be securely deployed."

[HWB14]





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 Nothing is disclosed throughout the computation



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- Can we **quantify** this leakage in a meaningful way?



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#### [Bac24, Part II]

- Develop an information-theoretic approach to quantify leakage
- Apply technique to a practically significant function(s)
- Determine and apply appropriate mitigation strategies

Function *f* evaluated on private data

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Partition parties P into:



Partition parties P into: attackers A,



Partition parties P into: attackers A, targets T,



Partition parties P into: attackers A, targets T, spectators S

Model participant inputs by a **random variable**  $X_{P_i}$ 

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C. Shannon. Photo: Alfred Eisenstaedt





#### Putting everything together

Attackers  $X_A$ , targets  $X_T$ , and spectators  $X_S$  (vectors)

Treat the **output** as a random variable:  $f(\mathbf{X}_A, \mathbf{X}_T, \mathbf{X}_S) = O$ 

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Attackers  $X_A$ , targets  $X_T$ , and spectators  $X_S$  (vectors) Treat the **output** as a random variable:  $f(X_A, X_T, X_S) = O$ 

# Attackers' weighted average entropy[AH17] $H(\mathbf{X}_T \mid \mathbf{X}_A = \mathbf{x}_A, O)$ how much information A learns<br/>about T, given $\mathbf{x}_A$ and OAbsolute entropy loss[BBZ24a; BBZ24b] $H(\mathbf{X}_T) - H(\mathbf{X}_T \mid \mathbf{X}_A = \mathbf{x}_A, O)$ the total amount of information<br/>disclosed about T, given $\mathbf{x}_A$ and O

Absolute entropy loss  $\iff$ mutual information between  $X_T$  and O (conditioned on  $X_A$ )

- 2016 Boston gender pay gap survey
- Analyzed the private wages based on gender and race using multi-party computation
- Average (salary) computation

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Source: Boston University, 2017

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$$O = \sum_{i} X_{T_i} + \sum_{j} X_{A_j} + \sum_{k} X_{S_k}$$

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**S**4





 $o = s_1 + s_2 + s_3 + s_4 + s_5$ 

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Are there certain input(s) an attacker can supply to **maximize** the information they learn?

#### Claim

The information disclosure is independent of the attackers' input(s).

- Intuition: an adversary can "remove" their influence
- Not universally true (depends on f)



- Model inputs by common distributions:
  - Poisson
  - Uniform
  - Gaussian
  - Log-normal [Cao+22]

- Model inputs by common distributions:
  - Poisson
  - Uniform
  - Gaussian
  - Log-normal [Cao+22]
- For a single evaluation, information disclosure is independent of
  - the distribution parameters
  - the distribution itself
- Disclosure is proportional to the number of **spectators**



Figure 1: Absolute entropy loss (lower is better), 1 target

- First study was a success
- Repeated the following year with an extended set of participants

The Boston Women's Workforce Council released its 2017

BOSTON WOMEN'S WORKFORCE COUNCIL REPORT 2017

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  - Correlated outputs

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#### An interesting question



#### An interesting question

"What happens if everyone else participates *again*, but without me?"



 $f(s_2, s_3, s_4, s_5) = o'$ 



Vary the **ratio** of shared spectators  $S_{1,2}$  to the (fixed) total number of spectators



Target's initial entropy
 After first evaluation
 Participating in both comps.
 Participating second comp. only
 Participating first comp. only

Figure 2: Conditional entropies, 6 total spectators, 1 target

Vary the **ratio** of shared spectators  $S_{1,2}$  to the (fixed) total number of spectators

- Largest protection at 50% overlap
- Undesirable disclosure at extrema
- Target's initial entropy
   After first evaluation
   Participating in both comps.
   Participating second comp. only
   Participating first comp. only





What are some logical successors to the average?

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- Order statistics (max/min, median)

$$f_{\max}(\mathbf{x}) = \max_i x_i$$

- Variability measures (variance)

$$f_{\sigma^2}(\mathbf{x}) = \frac{1}{n} \sum_{i} (x_i - f_{\mu}(\mathbf{x}))^2$$

- Multidimensional functions

$$f_{(\mu,\sigma^2)}(\mathbf{x}) = (f_{\mu}(\mathbf{x}), f_{\sigma^2}(\mathbf{x}))$$





#### Problem

Function could produce discrete outputs from continuous inputs, producing a "**mixture**"

## **Data-driven entropy estimators**



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Function could produce discrete outputs from continuous inputs, producing a "mixture"

#### Computer Science > Information Theory

[Submitted on 19 Sep 2017 (v1), last revised 9 Oct 2018 (this version, v3)]

#### Estimating Mutual Information for Discrete-Continuous Mixtures

Weihao Gao, Sreeram Kannan, Sewoong Oh, Pramod Viswanath

#### Data-driven entropy estimators



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Recall (from slide 9)...



An adversary **maximizes** the information learned by **minimizing** their influence.\*

<sup>\*</sup>Inverse behavior for the minimum

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In fact, the information A learns is **bounded** by observing the output, **without participating** in  $f_{max}$ 





Figure 3: Uniform  $\mathcal{U}(0,7)$ ,  $H(X_T | X_A = x_A, O)$ , 1 target

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Figure 3: Normal  $\mathcal{N}(0, 4)$ ,  $H(\mathbf{X}_T | \mathbf{X}_A = \mathbf{x}_A, O)$ , 1 target

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- Proof is a work-in-progress



Figure 3: Normal  $\mathcal{N}(0, 4)$ ,  $H(\mathbf{X}_T | \mathbf{X}_A = \mathbf{x}_A, O)$ , 1 target

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#### Further analysis of complex funcs.

- Derive analytical expressions information disclosure
- Apply data-driven analysis to broader functionalities

Data-driven analysis serves as a "first pass" assessment of a function's suitability for secure computation

Further analysis of complex funcs.	Mitigation strategies
<ul> <li>Derive analytical expressions information disclosure</li> </ul>	<ul> <li>Adding noise (differential privacy)</li> </ul>
<ul> <li>Apply data-driven analysis to broader functionalities</li> </ul>	<ul><li>Introducing synthetic inputs</li><li>Modifying the function</li></ul>

- Despite decades of performance improvements, broader privacy concerns remain that must be addressed prior to deployment of secure computation
- Developed a framework for quantifying information disclosure from secure computation outputs
- Computation designers can use this framework to determine potential disclosure about participants' inputs

# Thank you! Questions?

#### References

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